

Our work: • for all modalities of single-vector embeddings

- "theoretical limitation": "the need to transform to use multimodality" as the space of things to represent grows.

- limitations of vectors in geometric space:

~~History~~: ① The newly proposed angle:

Not How many k -subsets that a given configuration can realize

, but rather: what is the necessary embedding dimension to realize all k -subsets under score margin γ

- sphere-packing volume argument

↓
obtain a lower bound on the embedding dimension (in terms of n, k, γ)

Setup: Let $\{\vec{v}_i\}_{i=1}^n \subseteq \mathbb{R}^d$ be unit vectors. each representing a document.

Fix the score margin $\gamma > 0$.

We say that a k -subset $S \subseteq [n]$ is "realized with margin γ "

if $\exists$ unit query u_s , such that

$$\left(\min_{i \in S} \underbrace{\langle u_s, v_i \rangle}_{\leq \|u_s\| \cdot \|v_i\|} \right) \geq \left(\max_{j \notin S} \langle u_s, v_j \rangle \right) + 2 \times \gamma$$

$\|u_s\|, \|v_i\|$
both = 1

i.e., if the dot product between the specified query u_s and a candidate document (vector) is ~~the bigger~~ the better

then "realization of S " is existence of a "magic" u_s
s.t. lower "threshold" within S exceeds
upper "threshold" outside S by $\geq 2\gamma$

Thm 1: (lower bound of dimension d).

Assume $1 \leq k < n$. Assume $\forall$ k -subset $S \subseteq [n]$.

then S can be realized with margin γ .

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[this is an ideal assumption of document set $\{\vec{v}_i\}_{i=1}^n$]

Claim: $\binom{n}{k} \leq \left(1 + \frac{1}{\delta}\right)^d$, thus $d \leq \log_{\left(1 + \frac{1}{\delta}\right)} \binom{n}{k}$

proof: Consider $S, T \subseteq [n]$

where $|S| = |T| = K$

and $S \neq T$.

We can thus choose $(\exists) (i \in S \setminus T \text{ and } j \in T \setminus S)$.

*** In particular, by the assumption that every k -subset is γ -realizable, thus S and T are both γ -realizable

Thus $\exists$ unit query U_5 , where

$$\left(\min_{\tilde{i} \in S} \langle u_s, v_{\tilde{i}} \rangle \right) - \left(\max_{j \in S} \langle u_s, v_j \rangle \right) \geq 2\delta$$

and $\exists$ unit query $\boxed{u_I}$, where

$$\left(\min_{i \in T} \langle u_i, v_i \rangle \right) - \left(\max_{j \notin T} \langle u_j, v_j \rangle \right) \geq \delta$$

In particular, $\left\{ \begin{array}{l} \text{dot product with } u_S \text{ for some element in } S \\ \text{min in } S \\ \text{max not in } S \\ \text{dot prod. with } u_S \text{ for some element not in } S \end{array} \right\}$

$$\text{Thus } \begin{cases} \langle u_s, v_i \rangle - \langle u_s, v_j \rangle \geq 2 \times \gamma \\ \langle u_T, v_j \rangle - \langle u_T, v_i \rangle \geq 2 \times \gamma \end{cases}$$

$$\text{Thus } \begin{cases} \langle u_s, (v_i - v_j) \rangle \geq 2 \times \gamma \quad \text{--- ①} \\ \langle u_T, (v_j - v_i) \rangle \geq 2 \times \gamma \quad \text{--- ②} \end{cases}$$

adding ① and ②,

we have [whilst realizing $\langle u_T, (v_j - v_i) \rangle$
 $= \langle -u_T, (v_i - v_j) \rangle$

$\therefore$ we have $\langle (u_s - u_T), (v_i - v_j) \rangle$

$$\geq 2\gamma + 2\gamma = 4\gamma$$

Note that $4\gamma \leq \langle u_s - u_T, v_i - v_j \rangle$

we want to "say sth." about u_s and u_T

$$\begin{aligned} &\leq \|u_s - u_T\| \times \|v_i - v_j\| \\ &\leq \|u_s - u_T\| \times \left[\underbrace{\|v_i\|}_{=1} + \underbrace{\| -v_j \|}_{= \|v_j\| = 1} \right] \\ &\leq 2 \times \|u_s - u_T\| \end{aligned}$$

$$\therefore \|u_s - u_T\| \geq 2\gamma$$

because we assume all the n documents are represented as unit vectors)

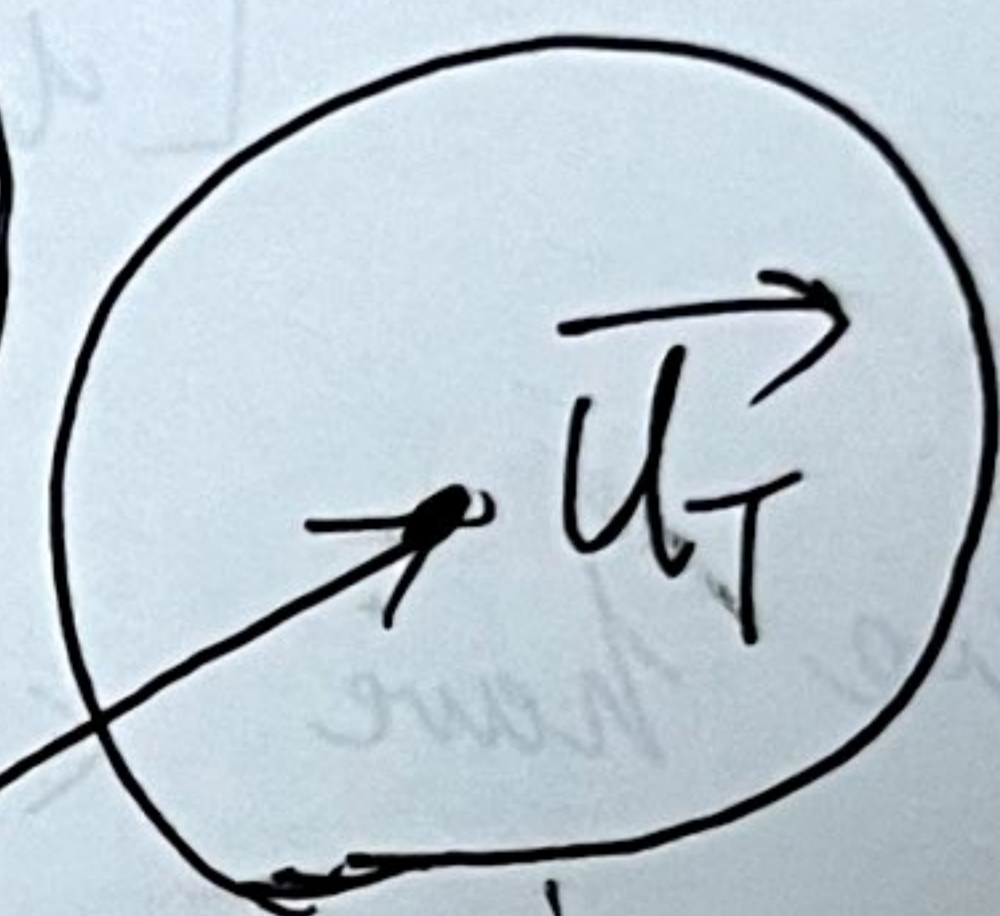
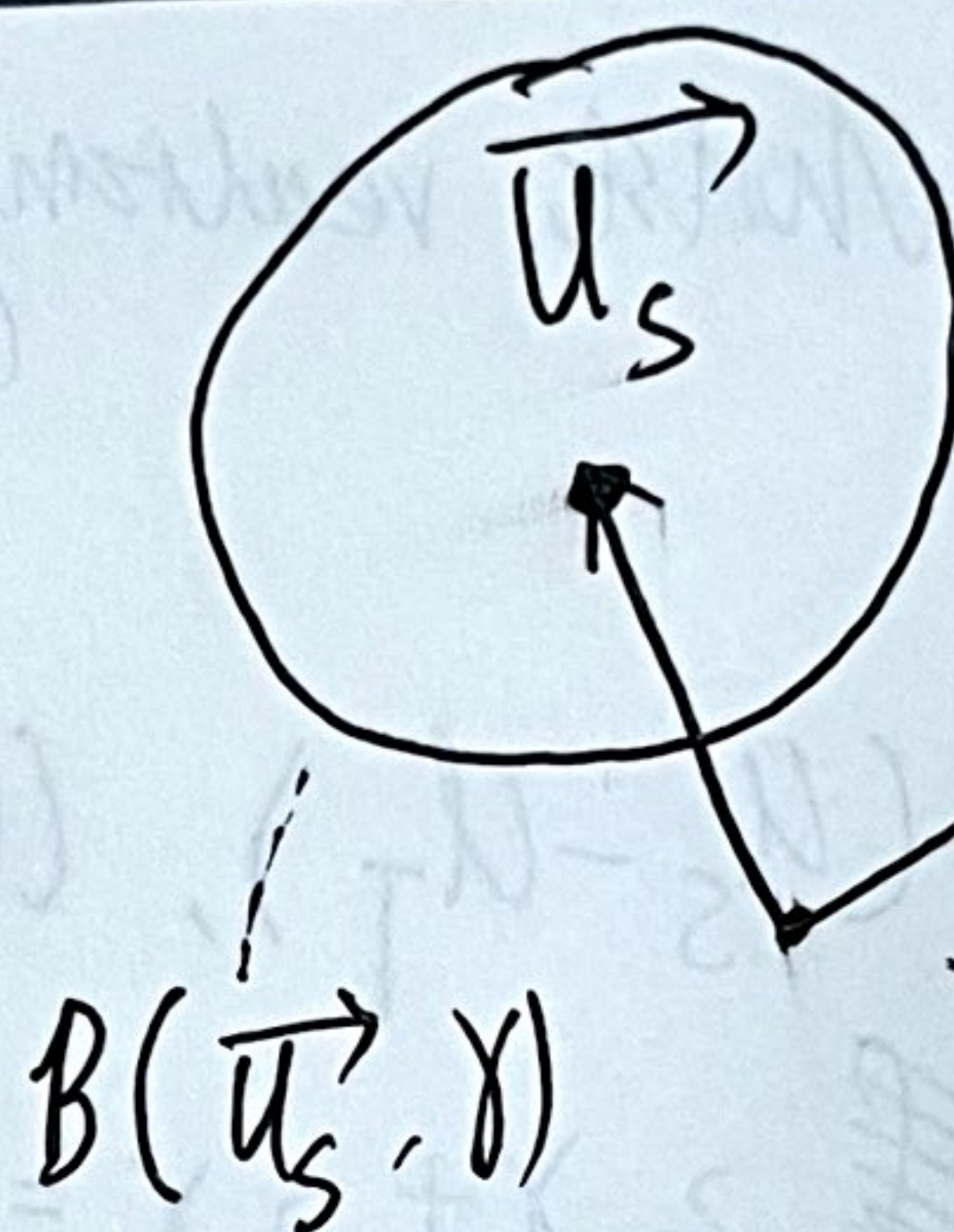
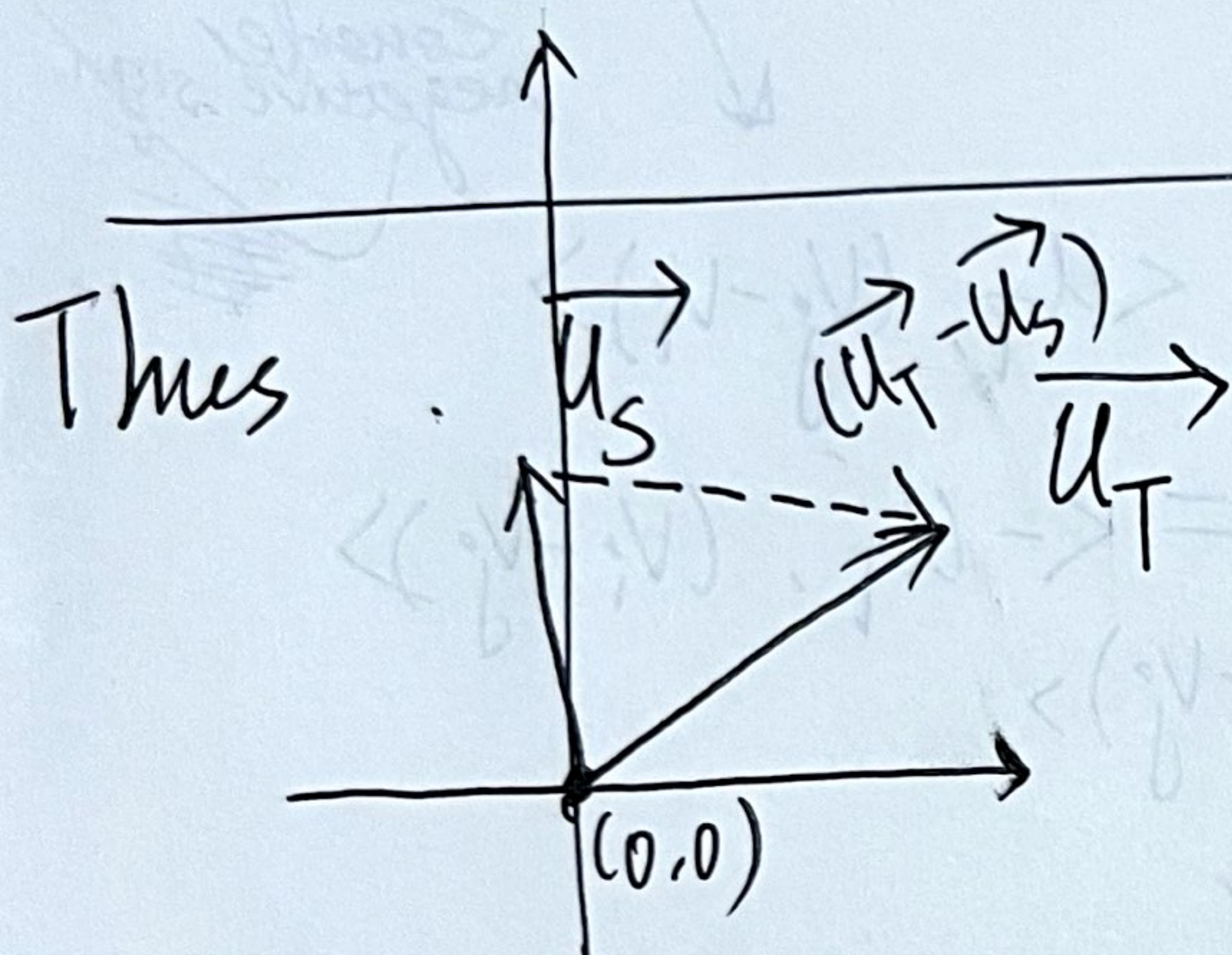
M_9 = number of size- K subsets S of $[n] := \{1, 2, \dots, n\}$
 $=$ # of queries u_s for such $S = \binom{n}{K}$

Recall that S and T are arbitrary size k subsets of $[n]$
 $(S \neq T)$

we know $\|u_S - u_T\| \geq 2 \times \gamma$

$\therefore$ In fact, for any 2 (out of M total)

k -subsets of S , their corresponding queries are operated by
under scheme of
 γ -realizability $\geq 2 \times \gamma$.



$\Rightarrow \|\vec{u}_T - \vec{u}_S\| \geq \gamma + \gamma$

center of ball
 $\vec{u}_S$

for $\forall S, T \subseteq [n]$
 $|S| = |T| = k$

Proposition: $B(u_S, \gamma) \subseteq B(0, (1+\gamma)\gamma)$

proof: suppose $\vec{x} \in B(u_S, \gamma)$
 be arbitrary.

Thus $\|\vec{x} - u_S\| \leq \gamma$

Thus $\|\vec{x} - \vec{0}\|$
 $= \|\vec{x}\| = \|(\vec{x} - u_S) + u_S\|$
 $\leq \|\vec{x} - u_S\| + \|u_S\|$
 $= 1 + \|\vec{x} - u_S\|$
 $\leq 1 + \gamma$

that the

open balls

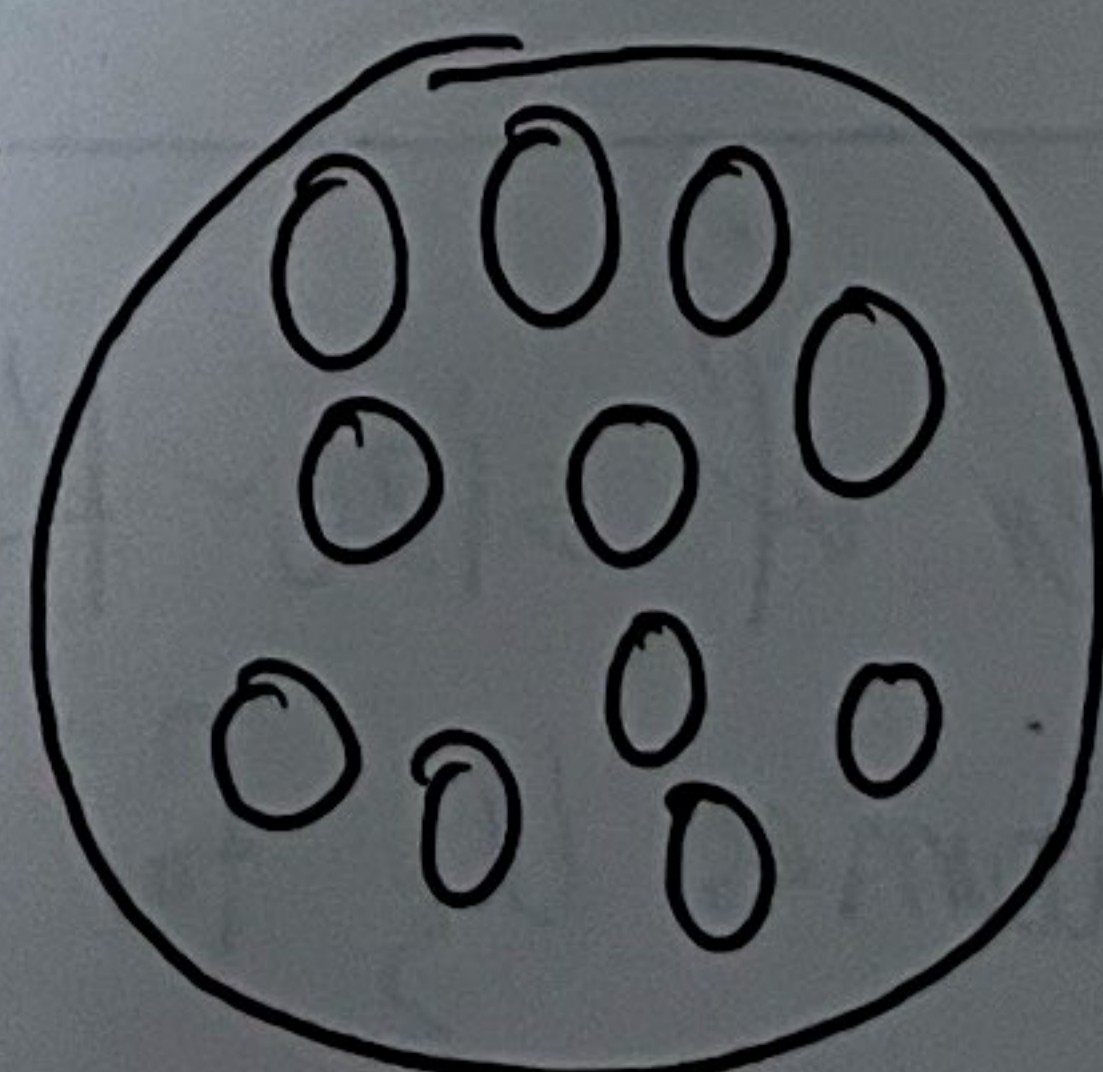
$B(\vec{u}_S, \gamma)$

and $B(\vec{u}_T, \gamma)$

are disjoint.

(i.e., $B(\vec{u}_S, \gamma) \cap B(\vec{u}_T, \gamma) = \emptyset$)

thus



$M \times V_d(\frac{\gamma}{2})$

$\leq V_d(1 + \gamma)$

$$\therefore M \times C_d \times \gamma^d \leq C_d \times (1+\gamma)^d$$

$$\therefore M \leq \left(\frac{\gamma}{1+\gamma}\right)^d = \left(1 - \frac{1}{1+\gamma}\right)^d$$